

DDoS and IP Traceback

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Denial Of Service

- ❑ The goal of a denial of service attack is to deny legitimate users access to a particular resource.
- ❑ An incident is considered an attack if a malicious user intentionally disrupts service to a computer or network resource.
- ❑ Resource exhaustion



- ❑ Distributed Denial of Service (DDoS)
- ❑ Proposed solutions
- ❑ Autonomous System (AS) based solution
- ❑ Conclusions

Resource Exhaustion

- ❑ Disk Space
- ❑ CPU Cycles
- ❑ Memory
- ❑ Network Bandwidth
- ❑ Application Resources
 - TCP Stack
 - Web Connections

Security

- ❑ People can justifiably rely on computer-based systems to perform critical functions
 - national scale infrastructures: water, power, communication, transportation, ...
 - localized systems: cars, homes, workplaces, ...
- ❑ People can justifiably rely on systems processing sensitive information about them to conform to public policy
 - health, banking, libraries, e-commerce, government records, ...
- ❑ *Without fear of sudden disruption by cyber attacks*

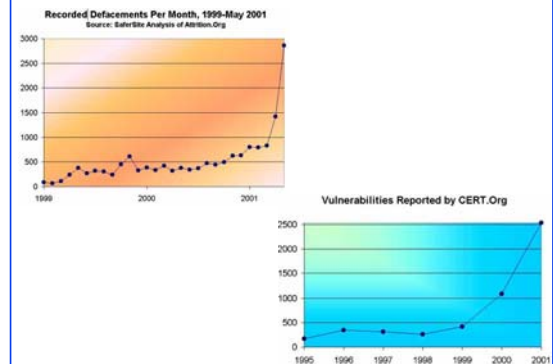
What's the Harm?

- ❑ Financial loss can be difficult to estimate
 - Lost business
 - Bad publicity and damaged reputation
- ❑ 2002 CSI/FBI Survey
 - 40% of reported attacks are DOS
 - Average cost per attack is >\$1 million
- ❑ Distributed DOS attacks (February 2000)
 - Amazon, CNN, E-Trade, eBay, etc...
 - Estimated losses were "several millions to billions of dollars"
- ❑ DOS can also be used to cover-up "real" attacks
- ❑ Nations critical infrastructure is also at risk

Denial of Service Attacks

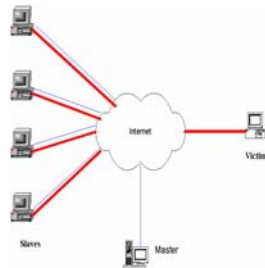
- ❑ Most involve either resource exhaustion or corruption of the operating system runtime environment.
- ❑ UDP bombing
- ❑ tcp SYN flooding
- ❑ ping of death
- ❑ smurf attack

Denial of Service



Distributed Denial of Service Attacks (DDoS)

- ❑ Attacker logs into Master and signals slaves to launch an attack on a specific target address (victim).
- ❑ Slaves then respond by initiating TCP, UDP, ICMP or Smurf attack on victim.



Why are DOS attacks possible?

- ❑ IP employs an open architecture
 - No authentication of sender's IP address
 - Easy to forge any address, hard to detect offender
 - IP traceback, ingress/egress filters (later)
- ❑ No resource regulation in the network
 - Employ QOS techniques to mitigate impact (later)

DDoS

- ❑ Denial of Service in pervasive networks
 - Power-draining attacks
 - Bandwidth-usage attacks
 - CPU-usage attack

Security Mechanisms

- ❑ Normally, not a single silver bullet
- ❑ Develop multiple layers of defense
- ❑ Employ as many layers of defense as needed - risk, resource profiles
- ❑ Castle, moat, drawbridge, mountain-top lookout, perimeter wall, inner wall, ruler decoy etc.
- ❑ Firewall, resource managers, app. Level security, logging, antivirus, remote backups, egress filters...

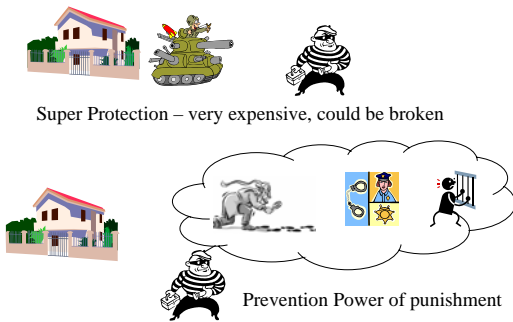
Security Analogy



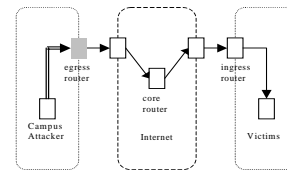
Resource Accounting

- ❑ Monitor network bandwidth, processor time and memory usage per process at server
- ❑ Regulate processes exceeding preset thresholds
- ❑ Problems: Hard to identify the process to whom resource usage needs to be charged
 - Interrupts, context-switches
 - A packet arrives at network interface

Two Security Philosophies



Ingress filtering



DOS attacks

- ❑ All DOS attacks consume resources
 - Bandwidth in UDP floods
 - Processing power in CGI bin attacks
 - Memory in fragmentation attacks
- ❑ Can we detect and contain attacks if we kept good accounting of resources?

DOS attacks

- ❑ Ingress filtering is not widely employed
 - Can be expensive in transit and backbone networks
- ❑ How to effectively trace back the source of the attack?
- ❑ If successful, may be able to throttle attack traffic at the network ingress

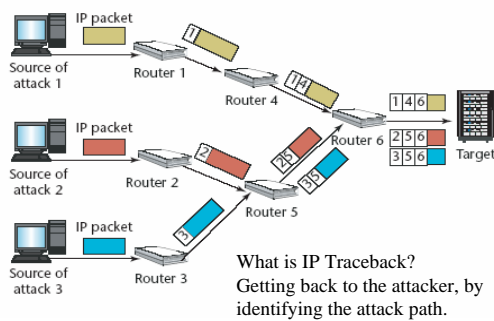
ICMP traceback(Bellovin, IETF)

- ❑ Generate ICMP packets with packet header, router and its neighbors ids
- ❑ Do this with low probability 1/20,000
- ❑ These ICMP packets can be used to trace the source
- ❑ More likely to get packets from routers closer to destination, rather than source

IP Traceback

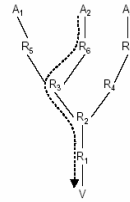
- ❑ Not practical to assume that all routers in the Internet will participate in marking scheme
- ❑ When some routers don't participate in marking, not sure if the last router in the constructed path is the true origin
- ❑ To be protected against single attacker that insert false information into the path the marking probability should be more than 0.5
 - Very high number (thousands) of packet to be analyzed by the victim

IP Traceback



IP traceback (Savage...Sigcomm00)

- ❑ Exact Traceback
 - R_6, R_3, R_2, R_1
- ❑ Approximate Traceback
 - Valid path suffix
 - R_5, R_6, R_3, R_2, R_1



IP Traceback

- ❑ Probabilistic Packet Marking (PPM)
 - No of attack packets required is 1000s
 - Difficult to handle DDoS attacks (to complex to construct attack path).
- ❑ ICMP Traceback or iTrace - Overhead
- ❑ Controlled Flooding - a form of DoS itself
- ❑ Hash-Based IP traceback
 - Less space needed and No eavesdropping
- ❑ IP Traceback with IPSec
 - Poor scalability
 - ISP need to update topology to all end users
 - End users need to know network topology

IP traceback -assumptions

- ❑ Attacker can generate any packet
- ❑ Attackers may conspire
- ❑ Aware of the tracing mechanism
- ❑ Attackers send lots of packets
- ❑ Packets may be lost, reordered
- ❑ Routes are pretty stable
- ❑ Routers are memory, CPU limited

IP traceback -Node Append

- ❑ Attach each router's IP address to the packet
 - Like IP record route option
- ❑ Every packet will have path info
- ❑ Too expensive
- ❑ Could lead to fragmentation problems

Edge Sampling

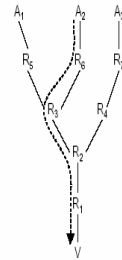
- ❑ Tree construction starting from victim (distance =0, 1,...)
- ❑ Time for convergence
 - furthest router: $p(1-p)^{d-1}$
- ❑ Can use any p, spoofed attacker packets distance field longer
- ❑ Robust against multiple attackers
 - Edges are different, linear complexity
- ❑ Takes many bits - $32+32+8? = 72$

Node Sampling

- ❑ Reserve a node field
- ❑ Routers write their IP address with probability p
- ❑ Prob. Of receiving id from d hops
 - $p(1-p)^{d-1}$
- ❑ $p > 0.5$, robust against attacker spoofing
- ❑ Routers far away from victim don't send many packets
 - $d=15, p=0.51$, expectation = 42,000 packets

Edge Sampling --encoding

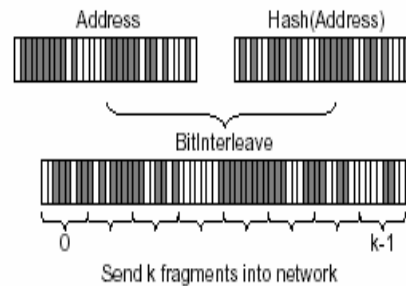
- ❑ Use XOR of addresses
- ❑ R1, 0
- ❑ R1 XOR R2, 1
- ❑ R1 XOR R2 XOR R3, 2
- ❑ Uses roughly half the space



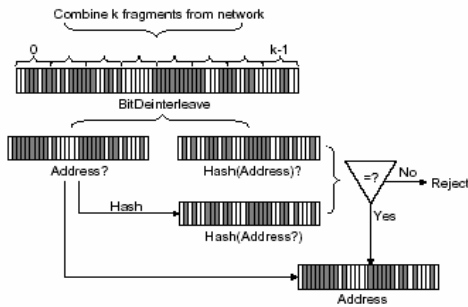
Edge Sampling

- ❑ Encode edges of path
 - Rather than single nodes
- ❑ Employ three fields
 - Start, end, distance
- ❑ With probability p, write Router IP address in start, make distance =0
- ❑ Else, (a) if start already marked, distance=0, put your id in end and
 - (b) increment distance

Edge Sampling— Fragment Sampling



Fragment Sampling



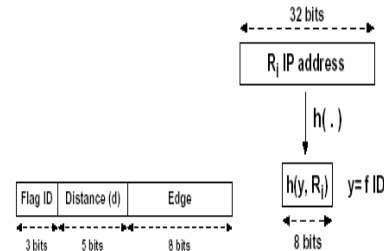
AMS

- Use two hash functions h and h'
- Encode $h(\text{start}) \text{ XOR } h'(\text{end})$
- Use 11-bits for hash, 5bits for length
- If you know upstream routers, few choices for $h(s)$, when we know $h'(e)$
- Tolerate multiple attackers
 - Upto 60
 - Main limitation: hash collisions

Fragment Sampling

- Can compress information into 16 bits
- Use IP fragment identifier space
- Expensive to compute
- Nor robust against large DDOS

AMS-II



Advanced Marking Scheme Song & Perrig, Infocom01

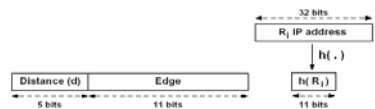
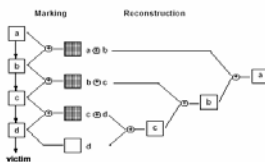


Figure 2: Encoding in Advanced Marking Scheme I



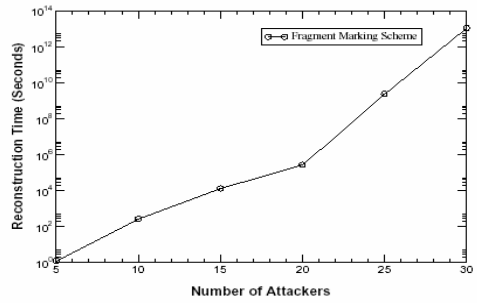
AMS-II

- Use two sets of hash functions
- Main intuition:
 - Probability of collision with 11 bits $1/2^{11}$
 - Probability of collision with m hashes of 11 bits = $1/(2^{11})^m$
 - Multiple hash functions reduce Collisions
- Where did we see that before?

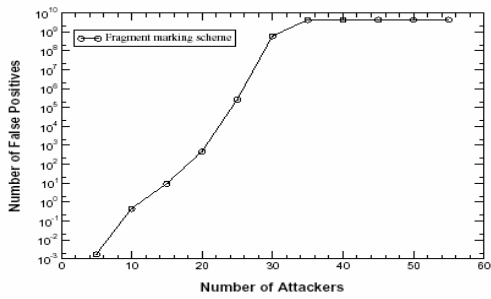
AMS-II

- ❑ Tries to work within the space of 11 bits
 - While identifying the hash function
- ❑ Easier than FSM
- ❑ Much more robust than FSM

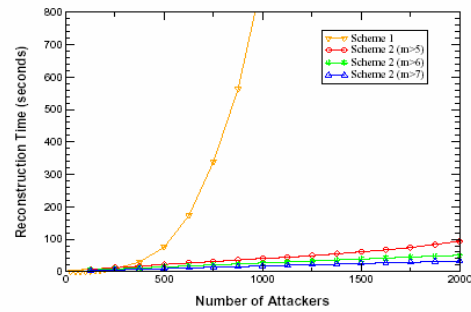
FMS Path reconstruction time



FMS False positives



AMS Path Reconstruction times



AMS & AMS-II

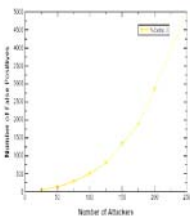


Figure 7: False Positives for Advanced Marking Scheme

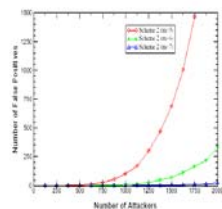
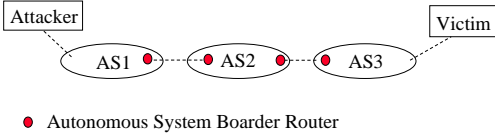


Figure 8: False Positives for Advanced Marking Scheme II

- ❑ Traceback is an interesting idea
- ❑ Allows us to trace the origin of the attack
- ❑ Threat of Identification leads to reduction in attacks
- ❑ What about the viruses?
 - Innocent attackers

Autonomous System - Traceback



Autonomous System Marking

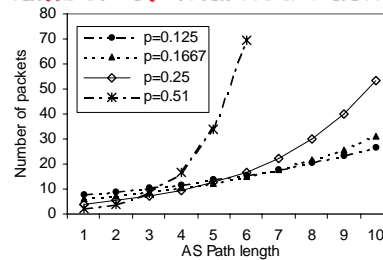
Marking procedure at router R with AS Number R_{AS} :

for each packet w
 let x be a random number from $[0, 1)$
 if $x < p$ then,
 write R_{AS} into $w.AS$
 set $w.AS_distance=0$
 else
 increment $w.AS_distance$

Autonomous Systems - AS

- AS is a group of IP networks managed by one network operator
- AS - set of routers using the same external routing policy
- Number of AS - 14,000, number of hosts - 200M
- In 99.5% of cases, a packet passes less than AS before reaching destination
- Network Operators may not always like to disclose their network details
- AS number is 16 bits compared to IP address 32 bits (IPv6 - 128 bits)

Number of Marked Packets



- $d_{AS} = 7, p=0.51 \Rightarrow 141$ packets needed
- If $p = 1/d_{AS} \Rightarrow 25$ packets needed

Autonomous System Marking

- Marking by ASBR
- Marking scheme similar to node sampling scheme
- 16 bits for ASN and 3 bits for AS_distance
- A packet is marked only if it leaves the AS
- A packet is marked with a probability p and the distance is set to zero
- If the ASBR chose not to mark, it increments the distance field

Authenticated Marking Scheme

- We assume the presence of a symmetric key infrastructure within each AS
- Each ASBR that belongs to the AS or connected to the AS know the secret key K_i
- Use one-way hash chains to generate session keys
 - h_0, h_1, \dots, h_n where $h_i = H(h_{i-1})$
 - Initially distribute h_0
 - Each ASBR computes the chain
 - Use the keys starting from the right to left

Authenticated AS Marking Algorithm

Marking procedure at router R with ASN R_{AS} :

K_{AS} is the symmetric key of R_{AS}

K'_{AS} is the symmetric key of the next AS in the path.

for each packet w

 Compute $D(AS\ Marking, K_{AS})$

 if (Redundancy Predicate is not fulfilled)

 Set AS Marking to $E(ASN \parallel RP, K'_{AS})$

 else

 let x be a random number from $[0, 1)$

 if $x < p$ then,

 Set AS marking to $E(ASN \parallel RP, K'_{AS})$

 else

 Set AS marking to $E(AS\ Marking, K'_{AS})$

Thank You!



Authenticated AS Traceback

- Victim obtains the AS symmetric key of the current session and computes AS marking
- Victim can reconstruct the path
- Victim can use the symmetric key to compute the keys of previous sessions but not any future sessions
 - A compromised victim doesn't affect the security of the mechanism

Summary



- Presented two schemes:
 - Autonomous System based Traceback
 - Authenticated Marking Scheme
- Only ASBR participate in marking
- Low marking overhead
- Enables to reconstruct the AS attack graph in real time
- Authenticated scheme prevents compromised routers from forging ASBR marking